The Total Irregularity of Graphs under Graph Operations

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Hosam Abdo, Darko Dimitrov

Institut für Informatik, Freie Universität Berlin, Takustraße 9, D-14195 Berlin, Germany E-mail: [abdo,darko]@mi.fu-berlin.de

Abstract

The total irregularity of a graph G is defined as $\operatorname{irr}_t(G) = \frac{1}{2} \sum_{u,v \in V(G)} |d_G(u) - d_G(v)|$, where $d_G(u)$ denotes the degree of a vertex $u \in V(G)$. In this paper we give (sharp) upper bounds on the total irregularity of graphs under several graph operations including join, lexicographic product, Cartesian product, strong product, direct product, corona product, disjunction and symmetric difference.

Keywords: Irregularity and total irregularity of graphs, graph operations

1 Introduction

Let G be a simple undirected graph with |V(G)| = n vertices and |E(G)| = m edges. The degree of a vertex v in G is the number of edges incident with v and it is denoted by $d_G(v)$. A graph G is regular if all its vertices have the same degree, otherwise it is irregular. However, in many applications and problems it is of big importance to know how irregular a given graph is. Several graph topological indices have been proposed for that purpose. Among the most investigated ones are: the irregularity of a graph introduced by Albertson [5], the variance of vertex degrees [8], and Collatz-Sinogowitz index [13].

The *imbalance* of an edge $e = uv \in E$, defined as $imb(e) = |d_G(u) - d_G(v)|$, appeares implicitly in the context of Ramsey problems with repeat degrees [6], and later in the work of Chen, Erdős, Rousseau, and Schlep [12], where 2-colorings of edges of a complete graph were considered. In [5], Albertson defined the *irregularity* of G as

$$irr(G) = \sum_{e \in E(G)} imb(e). \tag{1}$$

It is shown in [5] that for a graph G, $irr(G) < 4n^3/27$ and that this bound can be approached arbitrary closely. Albertson also presented upper bounds on irregularity for bipartite graphs,

triangle-free graphs and a sharp upper bound for trees. Some claims about bipartite graphs given in [5] have been formally proved in [19]. Related to Albertson is the work of Hansen and Mélot [18], who characterized the graphs with n vertices and m edges with maximal irregularity. The irregularity measure irr also is related to the first Zagreb index $M_1(G)$ and the second Zagreb index $M_2(G)$, one of the oldest and most investigated topological graph indices, defined as follows:

$$M_1(G) = \sum_{v \in V(G)} d_G^2(v)$$
 and $M_2(G) = \sum_{uv \in E(G)} d_G(u)d_G(v)$.

Alternatively the first Zagreb index can be expressed as

$$M_1(G) = \sum_{uv \in E(G)} [d_G(u) + d_G(v)].$$
 (2)

Fath-Tabar [15] established new bounds on the first and the second Zagreb indices that depend on the irregularity of graphs as defined in (1). In line with the standard terminology of chemical graph theory, and the obvious connection with the first and the second Zagreb indices, Fath-Tabar named the sum in (1) the third Zagreb index and denoted it by $M_3(G)$. The graphs with maximal irregularity with 6, 7 and 8 vertices are depicted in Figure 1.

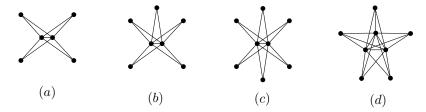


Figure 1: (a) The graph with 6 vertices with maximal irr. (b)The graph with 7 vertices with maximal irr. (c) and (d) Graphs with 8 vertices with maximal irr.

Two other most frequently used graph topological indices, that measure how irregular a graph is, are the variance of degrees and the Collatz-Sinogowitz index [13]. Let G be a graph with n vertices and m edges, and λ_1 be the index or largest eigenvalue of the adjacency matrix $A = (a_{ij})$ (with $a_{ij} = 1$ if vertices i and j are joined by an edge and 0 otherwise). Let n_i denotes the number of vertices of degree i for i = 1, 2, ..., n - 1. Then, the variance of degrees and the Collatz-Sinogowitz index are respectively defined as

$$\operatorname{Var}(G) = \frac{1}{n} \sum_{i=1}^{n-1} n_i \left(i - \frac{2m}{n} \right)^2 \quad \text{and} \quad \operatorname{CS}(G) = \lambda_1 - \frac{2m}{n}.$$
 (3)

Results of comparing irr, CS and Var are presented in [8, 14, 22].

There have been other attempts to determine how irregular graph is [2, 3, 4, 9, 10, 11, 23], but heretofore this has not been captured by a single parameter as it was done by the irregularity measure by Albertson.

The graph operation, especially graph products, plays significant role not only in pure and applied mathematics, but also in computer science. For example, the Cartesian product provide an important model for linking computers. In order to synchronize the work of the whole system it is necessary to search for Hamiltonian paths and cycles in the network. Thus, some results on Hamiltonian paths and cycles in Cartesian product of graphs can be applied in computer network design [27]. Many of the problems can be easily handled if the related graphs are regular or close to regular.

Recently in [1] a new measure of irregularity of a graph, so-called the *total irregularity*, that depends also on one single parameter (the pairwise difference of vertex degrees) was introduced. It was defined as

$$\operatorname{irr}_{t}(G) = \frac{1}{2} \sum_{u,v \in V(G)} |d_{G}(u) - d_{G}(v)|. \tag{4}$$

In the next theorem the upper bounds on the total irregularity of a graph are presented. Graphs with maximal total irregularity are depicted in Figure 2.

Theorem 1.1 ([1]). For a simple undirected graph G with n vertices, it holds that

$$\operatorname{irr}_t(G) \leq \begin{cases} \frac{1}{12}(2n^3 - 3n^2 - 2n) & n \text{ even,} \\ \\ \frac{1}{12}(2n^3 - 3n^2 - 2n + 3) & n \text{ odd.} \end{cases}$$

Moreover, the bounds are sharp.

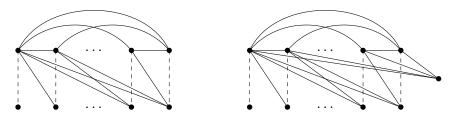


Figure 2: Graphs with maximal total irregularity H_n (with dashed edges) and \overline{H}_n (without dashed edges) for even and odd n, respectively.

The motivation to introduce the total irregularity of a graph, as modification of the irregularity of graph, is twofold. First, in contrast to irr(G), $irr_t(G)$ can be computed directly from the sequence of the vertex degrees (degree sequence) of G. Second, the most irregular graphs with respect to irr are graphs that have only two degrees (see Figure 1 for an illustration). On the contrary the most irregular graphs with respect to irr_t , as it is shown in [1], are graphs with maximal number of different vertex degrees (graphs with all doted (optional) edges in Figure 2), which is much closer to what one can expect from "very" irregular graphs.

The aim of this paper is to investigate the total irregularity of graphs under several graph operations including join, Cartesian product, direct product, strong product, lexicographic product, corona product, disjunction and symmetric difference. Detailed exposition on some graph operations one can find in [16].

2 Results

We start with simple observations about the complement and the disjoint union.

The complement of a simple graph G with n vertices, denoted by \overline{G} , is a simple graph with $V(\overline{G}) = V(G)$ and $E(\overline{G}) = \{uv \mid u, v \in V(G) \text{ and } uv \notin E(G)\}$. Thus, $uv \in E(G) \iff uv \notin E(G)$. Obviously, $E(G) \cup E(\overline{G}) = E(K_n)$, and for a vertex u, we have $d_{\overline{G}}(u) = n - 1 - d_G(u)$. From $|d_{\overline{G}}(u) - d_{\overline{G}}(v)| = |n - 1 - d_G(u) - (n - 1 - d_G(v))| = |d_G(u) - d_G(v)|$ it follows that $\operatorname{irr}_t(\overline{G}) = \operatorname{irr}_t(G)$.

For two graphs G_1 and G_2 with disjoint vertex sets $V(G_1)$ and $V(G_2)$ and disjoint edge sets $E(G_1)$ and $E(G_2)$ the disjoint union of G_1 and G_2 is the graph $G = G_1 \cup G_2$ with the vertex set $V(G_1) \cup V(G_2)$ and the edge set $E(G_1) \cup E(G_2)$. Obviously, $\operatorname{irr}_t(G \cup H) \geq \operatorname{irr}_t(G) + \operatorname{irr}_t(H)$.

Next we present sharp upper bounds for join, lexicographic product, Cartesian product, strong product, direct product, corona product and upper bounds for disjunction and symmetric difference.

2.1 Join

The join G + H of simple undirected graphs G and H is the graph with the vertex set $V(G + H) = V(G) \cup V(H)$ and the edge set $E(G + H) = E(G) \cup E(H) \cup \{uv : u \in V(G), v \in V(H)\}$.

Theorem 2.1. Let G and H be simple undirected graphs with $|V(G)| = n_1$ and $|V(H)| = n_2$ such that $n_1 \ge n_2$. Then,

$$\operatorname{irr}_{t}(G+H) \leq \operatorname{irr}_{t}(G) + \operatorname{irr}_{t}(H) + n_{2}(n_{1}-1)(n_{1}-2).$$

Moreover, the bound is best possible.

Proof. The total irregularity of G + H is

$$\operatorname{irr}_{t}(G+H) = \frac{1}{2} \sum_{u,v \in V(G+H)} |d_{G+H}(u) - d_{G+H}(v)|$$

$$= \frac{1}{2} \sum_{u,v \in V(G)} |d_{G+H}(u) - d_{G+H}(v)|$$

$$+ \frac{1}{2} \sum_{u,v \in V(H)} |d_{G+H}(u) - d_{G+H}(v)|$$

$$+ \sum_{u \in V(G)} \sum_{v \in V(H)} |d_{G+H}(u) - d_{G+H}(v)|.$$

By definition, $|V(G+H)| = |V(G)| + |V(H)| = n_1 + n_2$. For vertices $u \in V(G)$ and $v \in V(H)$, it holds that $d_{G+H}(u) = d_G(u) + n_2$ and $d_{G+H}(v) = d_H(v) + n_1$. Thus, further we have

$$\operatorname{irr}_{t}(G+H) = \frac{1}{2} \sum_{u,v \in V(G)} |d_{G}(u) - d_{G}(v)| + \frac{1}{2} \sum_{u,v \in V(H)} |d_{H}(u) - d_{H}(v)| + \sum_{u \in V(G)} \sum_{v \in V(H)} |(d_{G}(u) + n_{2}) - (d_{H}(v) + n_{1})|$$
(5)

$$= \operatorname{irr}_{t}(G) + \operatorname{irr}_{t}(H) + \sum_{u \in V(G)} \sum_{v \in V(H)} |n_{1} - n_{2} + d_{H}(v) - d_{G}(u)|.$$

Under the constrains $n_1 \geq n_2$, $d_G(u) \leq n_1 - 1$, and $d_H(v) \leq n_2 - 1$, the double sum $\sum_{u \in V(G)} \sum_{v \in V(H)} |n_1 - n_2 + d_H(v) - d_G(u)|$ is maximal when H is a graph with maximal sum of vertex degrees, i.e., H is the complete graph K_{n_2} , and G is a graph with minimal sum of

vertex degrees, i.e., G is a tree on n_1 vertices T_{n_1} . Thus,

$$\sum_{u \in V(G)} \sum_{v \in V(H)} |n_1 - n_2 + d_H(v) - d_G(u)|$$

$$\leq \sum_{u \in V(T_{n_1})} \sum_{v \in V(K_{n_2})} |n_1 - n_2 + d_{K_{n_2}}(v) - d_{T_{n_1}}(u)|$$

$$= \sum_{u \in V(T_{n_1})} \sum_{v \in V(K_{n_2})} |n_1 - 1 - d_{T_{n_1}}(u)|$$

$$= n_2 \sum_{u \in V(T_{n_1})} (n_1 - 1 - d_{T_{n_1}}(u))$$

$$= n_2 n_1 (n_1 - 1) - 2n_2 (n_1 - 1)$$

$$= n_2 (n_1 - 1) (n_1 - 2).$$

and

$$\operatorname{irr}_t(G+H) \le \operatorname{irr}_t(G) + \operatorname{irr}_t(H) + n_2(n_1-1)(n_1-2).$$
 (6)

When $n_1 \leq 2$, $\operatorname{irr}_t(G) = \operatorname{irr}_t(H) = \operatorname{irr}_t(G+H) = 0$, and the claim of the theorem is fulfilled. From the derivation, it follows that (6) is equality when H is compete graph on n_2 vertices and G is any tree on n_1 vertices.

Example. Let denote by H_i a graph with $|V(H_i)| = i$ isolated vertices (vertices with degree zero). Then, the bipartite graph $K_{i,j}$ is a join of H_i and H_j . Analogously, the complete k-partite graph $G = K_{n_1,\dots,n_k}$ is join of H_{n_1},\dots,H_{n_k} . Straightforward calculation shows that $\operatorname{irr}_t(K_{n_i,n_j}) = n_i n_j |n_j - n_i|$, For the total irregularity of K_{n_1,\dots,n_k} we have

$$\operatorname{irr}_{t}(K_{n_{1},\dots,n_{k}}) = \frac{1}{2} \sum_{u,v \in V(K_{n_{1},\dots,n_{k}})} |d_{G}(u) - d_{G}(v)|$$

$$= \sum_{i=1}^{k-1} \sum_{j=i+1}^{k} \left(\frac{1}{2} \sum_{u,v \in V(K_{n_{i},n_{j}})} |d_{G}(u) - d_{G}(v)| \right)$$

$$= \sum_{i=1}^{k-1} \sum_{j=i+1}^{k} n_{i} n_{j} |n_{j} - n_{i}| = \sum_{i=1}^{k-1} \sum_{j=i+1}^{k} \operatorname{irr}_{t}(K_{n_{i},n_{j}}).$$

2.2 Lexicographic product

The lexicographic product $G \circ H$ (also known as the graph composition) of simple undirected graphs G and H is the graph with the vertex set $V(G \circ H) = V(G) \times V(H)$ and the edge set $E(G \circ H) = \{(u_i, v_k)(u_j, v_l) : [u_i u_j \in E(G)] \vee [(v_k v_l \in E(H)) \wedge (u_i = u_j)]\}.$

Theorem 2.2. Let G and H be simple undirected graphs with $|V(G)| = n_1$, $|V(H)| = n_2$ then,

$$\operatorname{irr}_t(G \circ H) \leq n_2^3 \operatorname{irr}_t(G) + n_1^2 \operatorname{irr}_t(H).$$

Moreover, this bound is sharp for infinitely many graphs.

Proof. By the definition of $G \circ H$, it follows that $|V(G \circ H)| = n_1 n_2$ and $d_{G \circ H}(u_i, v_j) = n_2 d_G(u_i) + d_H(v_j)$ for all $1 \le i \le n_1, 1 \le j \le n_2$. Applying those relations, we obtain

$$\operatorname{irr}_{t}(G \circ H) = \frac{1}{2} \sum_{\substack{(u_{i}, v_{k}) \in V(G \circ H) \\ (u_{j}, v_{l}) \in V(G \circ H)}} |d_{G \circ H}(u_{i}, v_{k}) - d_{G \circ H}(u_{j}, v_{l})|$$

$$= \frac{1}{2} \sum_{\substack{u_{i}, u_{j} \in V(G) \\ v_{k}, v_{l} \in V(H)}} |n_{2} d_{G}(u_{i}) - n_{2} d_{G}(u_{j}) + d_{H}(v_{k}) - d_{H}(v_{l})|$$

$$\leq \frac{1}{2} \sum_{\substack{u_{i}, u_{j} \in V(G) \\ v_{k}, v_{l} \in V(H)}} (n_{2} |d_{G}(u_{i}) - d_{G}(u_{j})| + |d_{H}(v_{k}) - d_{H}(v_{l})|)$$

$$= \frac{1}{2} n_{2}^{3} \sum_{\substack{u_{i}, u_{j} \in V(G) \\ v_{k}, v_{l} \in V(H)}} |d_{G}(u_{i}) - d_{G}(u_{j})|$$

$$+ \frac{1}{2} n_{1}^{2} \sum_{\substack{v_{k}, v_{l} \in V(H) \\ v_{k}, v_{l} \in V(H)}} |d_{H}(v_{k}) - d_{H}(v_{l})|$$

$$= n_{2}^{3} \operatorname{irr}_{t}(G) + n_{1}^{2} \operatorname{irr}_{t}(H). \tag{7}$$

To prove that the presented bound is best possible, consider the lexicographic product $P_l \circ C_k$, $l \geq 1, k \geq 3$ (an illustration is given in Figure 3(b)). Straightforward calculations give that $\operatorname{irr}_t(P_l) = 2(l-2)$, $\operatorname{irr}_t(C_k) = 0$. The graph $P_l \circ C_k$ is comprised of 2k vertices of degree k+2, and k(l-2) vertices of degree 2k+2. Hence, $\operatorname{irr}_t(P_l \circ C_k) = 2k^3(l-2)$. On the other hand, the bound obtain here is $\operatorname{irr}_t(P_l \circ C_k) \leq k^3 \operatorname{irr}_t(P_l) + l^2 \operatorname{irr}_t(C_k) = 2k^3(l-2)$.

2.3 Cartesian product

The Cartesian product $G \square H$ of two simple undirected graphs G and H is the graph with the vertex set $V(G \square H) = V(G) \times V(H)$ and the edge set $E(G \square H) = \{(u_i, v_k)(u_j, v_l) : [(u_i u_j \in E(G)) \land (v_k = v_l)] \lor [(v_k v_l \in E(H)) \land (u_i = u_j)] \}$. From the definition of the Cartesian product, it follows that $|V(G \square H)| = n_1 n_2$ and $d_{G \square H}(u_i, v_j) = d_G(u_i) + d_H(v_j)$. Since the derivation of the upper bound on $G \square H$ is similar to the case of a graph lexicographic product, we omit the proof and just state the result in Theorem 2.3. The best possible bound is obtained for $P_l \square C_k$, $l \ge 1, k \ge 3$, illustrated in Figure 3(c). The graph $P_l \square C_k$ is comprised of 2k vertices of degree 3, and k(l-2) vertices of degree 4. Thus, $\operatorname{irr}_t(P_l \square C_k) = 2k^2(l-2)$. The bound obtain here is $\operatorname{irr}_t(P_l \square C_k) \le k^2 \operatorname{irr}_t(P_l) + l^2 \operatorname{irr}_t(C_k) = 2k^2(l-2)$.

Theorem 2.3. Let G and H be simple undirected graphs with $|V(G)| = n_1$, $|V(H)| = n_2$ then

$$\operatorname{irr}_t(G \square H) \le n_2^2 \operatorname{irr}_t(G) + n_1^2 \operatorname{irr}_t(H).$$

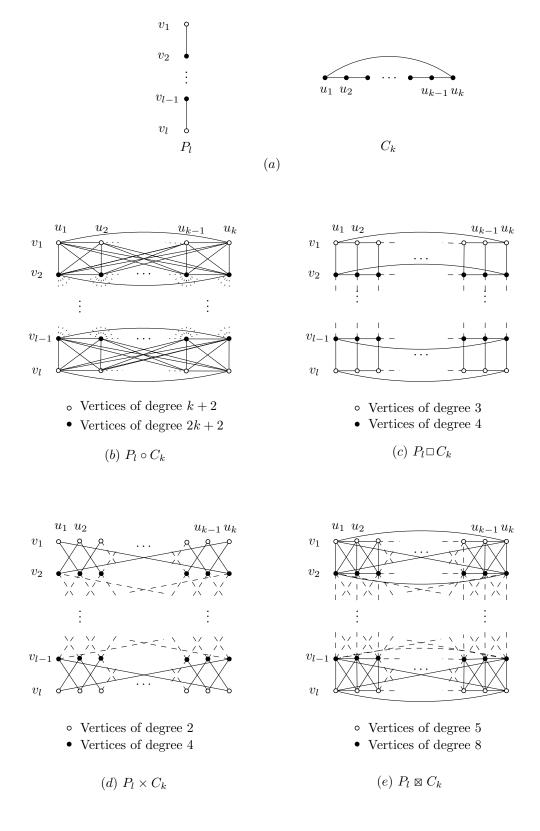


Figure 3: (a) Path graph on l vertices P_l , and cycle graph on k vertices C_k , (b) lexicographic product graph $P_l \circ C_k$, (c) Cartesian product graph $P_l \square C_k$, (d) direct product graph $P_l \times C_k$ and, (e) strong product graph $P_l \boxtimes C_k$.

Moreover, this bound is sharp for infinitely many graphs.

2.4 Strong product

The strong product $G \boxtimes H$ of two simple undirected graphs G and H is the graph with the vertex set $V(G \boxtimes H) = V(G) \times V(H)$ and the edge set $E(G \boxtimes H) = \{(u_i, v_k)(u_j, v_l) : [(u_i u_j \in E(G)) \land (v_k = v_l)] \lor [(v_k v_l \in E(H)) \land (u_i = u_j)] \lor [(u_i u_j \in E(G)) \land (v_k v_l \in E(H))] \}.$

Theorem 2.4. Let G and H be simple undirected graphs with $|V(G)| = n_1$, $|E(G)| = m_1$, $|V(H)| = n_2$ and $|E(H)| = m_2$. Then,

$$\operatorname{irr}_t(G \boxtimes H) \leq n_2 (n_2 + 2m_2) \operatorname{irr}_t(G) + n_1 (n_1 + 2m_1) \operatorname{irr}_t(H).$$

Moreover, this bound is best possible.

Proof. From the definition of the strong product, it follows $|V(G \boxtimes H)| = n_1 n_2$, $|E(G \boxtimes H)| = m_1 n_2 + m_2 n_1 + 2m_1 m_2$, and $d_{G \boxtimes H}(u_i, v_k) = d_G(u_i) + d_H(v_k) + d_G(u_i) d_H(v_k)$. The total irregularity of $G \boxtimes H$ is

$$\operatorname{irr}_{t}(G \boxtimes H) = \frac{1}{2} \sum_{(u_{i}, v_{k}), (u_{j}, v_{l}) \in V(G \boxtimes H)} |d_{G \boxtimes H}(u_{i}, v_{k}) - d_{G \boxtimes H}(u_{j}, v_{l})|$$

$$= \frac{1}{2} \sum_{u_{i}, u_{j} \in V(G), v_{k}, v_{l} \in V(H)} |d_{G \boxtimes H}(u_{i}, v_{k}) - d_{G \boxtimes H}(u_{j}, v_{l})|. \tag{8}$$

Applying simple algebraic transformation and the triangle inequality, we obtain

$$|d_{G\boxtimes H}(u_{i}, v_{k}) - d_{G\boxtimes H}(u_{j}, v_{l})| = |(d_{G}(u_{i}) - d_{G}(u_{j})) + (d_{H}(v_{k}) - d_{H}(v_{l})) + (d_{G}(u_{i})d_{H}(v_{k}) - d_{G}(u_{j})d_{H}(v_{l}))|$$

$$\leq |d_{G}(u_{i}) - d_{G}(u_{j})| + |d_{H}(v_{k}) - d_{H}(v_{l})|$$

$$+ \frac{1}{2}(d_{G}(u_{i}) + d_{G}(u_{j})) |d_{H}(v_{k}) - d_{H}(v_{l})|$$

$$+ \frac{1}{2}(d_{H}(v_{k}) + d_{H}(v_{l})) |d_{G}(u_{i}) - d_{G}(u_{j})|.$$
 (9)

From (8) and (9), we obtain

$$\operatorname{irr}_{t}(G \boxtimes H) \leq \frac{1}{2} \sum_{\substack{u_{i}, u_{j} \in V(G), \\ v_{k}, v_{l} \in V(H)}} [|d_{G}(u_{i}) - d_{G}(u_{j})| + |d_{H}(v_{l}) - d_{H}(v_{k})|]$$

$$+ \frac{1}{4} \sum_{\substack{u_{i}, u_{j} \in V(G), \\ v_{k}, v_{l} \in V(H)}} (d_{G}(u_{i}) + d_{G}(u_{j})) |d_{H}(v_{k}) - d_{H}(v_{l})|$$

$$+ \frac{1}{4} \sum_{\substack{u_{i}, u_{j} \in V(G), \\ v_{k}, v_{l} \in V(H)}} (d_{H}(v_{k}) + d_{H}(v_{l})) |d_{G}(u_{i}) - d_{G}(u_{j})|$$

$$= n_{2} (n_{2} + 2m_{2}) \operatorname{irr}_{t}(G) + n_{1} (n_{1} + 2m_{1}) \operatorname{irr}_{t}(H).$$

To prove that the presented bound is best possible, consider the strong product $P_l \otimes C_k$, $l \geq 1, k \geq 3$, illustrated in Figure 3(e). We have, $\operatorname{irr}_t(P_l) = 2(l-2)$, $\operatorname{irr}_t(C_k) = 0$. The graph $P_l \boxtimes C_k$ is comprised of 2k vertices of degree 5, and k(l-2) vertices of degree 8. Hence, $\operatorname{irr}_t(P_l \boxtimes C_k) = 6k^2(l-2)$. On the other hand, the bound obtain here, is $\operatorname{irr}_t(P_l \boxtimes C_k) \leq l(l+2(l-1))\operatorname{irr}_t(C_k) + k(k+2k)\operatorname{irr}_t(P_l) = 6k^2(l-2)$.

2.5 Direct product

The direct product $G \times H$ (also know as the tensor product, the Kronecker product [26], categorical product [25] and conjunctive product) of simple undirected graphs G and H is the graph with the vertex set $V(G \times H) = V(G) \times V(H)$, and the edge set $E(G \times H) = \{(u_i, v_k)(u_j, v_l) : (u_i, u_j) \in E(G) \land (v_k, v_l) \in E(H)\}$. From the definition of the direct product, it follows $|V(G \times H)| = n_1 n_2$, $|E(G \times H)| = 2m_1 m_2$, and $d_{G \times H}(u_i, v_k) = d_G(u_i) d_H(v_k)$. The proof for the upper bound on $G \times H$ is similar as that of the strong product $G \boxtimes H$. Therefore, we show only that the bound in Theorem 2.5 is best possible, and omit the rest of the proof.

Theorem 2.5. Let G and H be simple undirected graphs with $|V(G)| = n_1$, $|E(G)| = m_1$, $|V(H)| = n_2$ and $|E(H)| = m_2$. Then,

$$\operatorname{irr}_t(G \times H) \le 2 n_2 m_2 \operatorname{irr}_t(G) + 2 n_1 m_1 \operatorname{irr}_t(H).$$

Moreover, this bound is best possible.

To prove that the presented bound is best possible, we consider the direct product $P_l \times C_k$, $l \ge 1, k \ge 3$ (an illustration is given in Figure 3(d). Straightforward calculations give that $\operatorname{irr}_t(P_l) = 2(l-2)$, $\operatorname{irr}_t(C_k) = 0$. The graph $P_l \times C_k$ is comprised of 2k vertices of degree 2, and k(l-2) vertices of degree 4. Thus, $\operatorname{irr}_t(P_l \times C_k) = 4k^2(l-2)$. On the other hand, the bound obtain by Proposition 2.5 is $\operatorname{irr}_t(P_l \times C_k) \le 2n_2m_2\operatorname{irr}_t(G) + 2n_1m_1\operatorname{irr}_t(H) = 4k^2(l-2)$.

2.6 Corona product

The corona product $G \odot H$ of simple undirected graphs G and H with $|V(G)| = n_1$ and $|V(H)| = n_2$, is defined as the graph who is obtained by taking the disjoint union of G and n_1 copies of H and for each i, $1 \le i \le n_1$, inserting edges between the ith vertex of G and each vertex of the ith copy of G. Thus, the corona graph $G \odot H$ is the graph with the vertex set $V(G \odot H) = V(G) \cup_{i=1,\dots,n_1} V(H_i)$ and the edge set $E(G \odot H) = E(G) \cup_{i=1,\dots,n_1} E(H_i) \cup \{u_i v_j : u_i \in V(G), v_j \in V(H_i)\}$, where H_i is the ith copy of the graph H.

Theorem 2.6. Let G and H be simple undirected graphs with $|V(G)| = n_1$ and $|V(H)| = n_2$. Then,

$$\operatorname{irr}_t(G \odot H) \le \operatorname{irr}_t(G) + n_1^2 \operatorname{irr}_t(H) + n_1^2 (n_2^2 + n_1 n_2 - 4n_2 + 2).$$

Moreover, the bound is best possible.

Proof. The total irregularity of $G \odot H$ is

$$\operatorname{irr}_{t}(G \odot H) = \frac{1}{2} \sum_{u,v \in V(G \odot H)} |d_{G \odot H}(u) - d_{G \odot H}(v)|$$

$$= \frac{1}{2} \sum_{x,y \in V(G)} |d_{G \odot H}(x) - d_{G \odot H}(y)|$$

$$+ \sum_{i=1}^{n_{1}} \left(\frac{1}{2} \sum_{z,t \in V(H_{i})} |d_{G \odot H}(z) - d_{G \odot H}(t)| \right)$$

$$+ \sum_{i=1}^{n_{1}-1} \sum_{j=i+1}^{n_{1}} \sum_{z \in V(H_{i}),t \in V(H_{j})} |d_{G \odot H}(z) - d_{G \odot H}(t)|$$

$$+ \sum_{i=1}^{n_{1}} \sum_{u \in V(G),v \in V(H)} |d_{G \odot H}(u) - d_{G \odot H}(v)|.$$

By the definition of $G \odot H$, $|V(G \odot H)| = |V(G)| + n_1 |V(H)| = n_1 + n_1 n_2$. For a vetrex $u \in V(G)$, it holds that $d_{G \odot H}(u) = d_G(u) + n_2$ and for a vertex $v \in V(H_i)$, $1 \le i \le n_2$, we have $d_{G \odot H}(v) = d_H(v) + 1$. Thus,

$$\operatorname{irr}_{t}(G \odot H) = \frac{1}{2} \sum_{x,y \in V(G)} |d_{G}(x) - d_{G}(y)| + \frac{1}{2} n_{1} \sum_{z,t \in V(H)} |d_{H}(z) - d_{H}(t)|$$

$$+ \sum_{i=1}^{n_{1}-1} \sum_{j=i+1}^{n_{1}} \sum_{z \in V(H_{i}), t \in V(H_{j})} |d_{H}(z) + 1 - d_{H}(t) - 1|$$

$$+ \sum_{i=1}^{n_{1}} \sum_{u \in V(G), v \in V(H)} |d_{G}(u) - d_{H}(v) + n_{2} - 1|$$

$$= \operatorname{irr}_{t}(G) + n_{1} \operatorname{irr}_{t}(H) + n_{1}(n_{1} - 1) \operatorname{irr}_{t}(H)$$

$$+ \sum_{i=1}^{n_{1}} \sum_{u \in V(G), v \in V(H)} |d_{G}(u) - d_{H}(v) + n_{2} - 1|.$$

$$(10)$$

Since $n_1 \ge n_2$, the sum $\sum_{u \in V(G), v \in V(H)} |d_G(u) - d_H(v) + n_2 - 1|$ is maximal when $\sum_{u \in V(G)} d_G(u)$ is maximal, i.e., G is the complete graph K_{n_1} , and $\sum_{v \in V(H)} d_H(v)$ is minimal, i.e., H is a tree on n_2 vertices T_{n_2} . Thus,

$$\sum_{u \in V(G), v \in V(H)} |d_G(u) - d_H(v) + n_2 - 1|$$

$$\leq \sum_{u \in V(K_{n_1})} \sum_{v \in V(T_{n_2})} |d_{K_{n_1}}(u) - d_{T_{n_2}}(v) + n_2 - 1|$$

$$= \sum_{u \in V(K_{n_1})} \sum_{v \in V(T_{n_2})} |n_1 - 1 - d_{T_{n_2}}(v) + n_2 - 1|$$

$$= n_1 \sum_{v \in V(T_{n_2})} (n_1 + n_2 - 2 - d_{T_{n_2}}(v))$$

$$= n_1 n_2 (n_1 + n_2 - 2) - 2n_1 (n_2 - 1)$$

$$= n_1 (n_2^2 + n_1 n_2 - 4n_2 + 2). \tag{11}$$

Substituting (11) into (10), we obtain

$$\operatorname{irr}_t(G \odot H) \le \operatorname{irr}_t(G) + n_1^2 \operatorname{irr}_t(H) + n_1^2 \left(n_2^2 + n_1 n_2 - 4n_2 + 2\right).$$
 (12)

From the derivation of the bound (12), it follows that the sharp bound is obtained when G is compete graph on n_1 vertices and H is any tree on n_2 vertices.

2.7 Disjunction

The disjunction graph $G \vee H$ of simple undirected graphs G and H with $|V(G)| = n_1$ and $|V(H)| = n_2$ is the graph with the vertex set $V(G \vee H) = V(G) \times V(H)$ and the edge set $E(G \vee H) = \{(u_i, v_k)(u_j, v_l) : u_i u_j \in E(G) \vee v_k v_l \in E(H)\}$. It holds that $|V(G \vee H)| = n_1 n_2$, and $d_{G \vee H}(u_i, v_k) = n_2 d_G(u_i) + n_1 d_H(v_k) - d_G(u_i) d_H(v_k)$ for all i, k where, $1 \leq i \leq n_1$, $1 \leq k \leq n_2$.

Theorem 2.7. Let G and H be simple undirected graphs with $|V(G)| = n_1$, $|E(G)| = m_1$, $|V(H)| = n_2$ and $|E(H)| = m_2$. Then,

$$\operatorname{irr}_t(G \vee H) \leq n_2 (n_2^2 + 2m_2) \operatorname{irr}_t(G) + n_1 (n_1^2 + 2m_1) \operatorname{irr}_t(H).$$

Proof. The total irregularity of $G \vee H$ is

$$\operatorname{irr}_{t}(G \vee H) = \frac{1}{2} \sum_{(u_{i}, v_{k}), (u_{j}, v_{l}) \in V(G \vee H)} |d_{G \vee H}(u_{i}, v_{k}) - d_{G \vee H}(u_{j}, v_{l})|$$

$$= \frac{1}{2} \sum_{u_{i}, u_{j} \in V(G), v_{k}, v_{l} \in V(H)} |d_{G \vee H}(u_{i}, v_{k}) - d_{G \vee H}(u_{j}, v_{l})|. \tag{13}$$

Since $d_{G\vee H}(u_i, v_k) = n_2 d_G(u_i) + n_1 d_H(v_k) - d_G(u_i) d_H(v_k)$ for all i, k where, $1 \leq i \leq n_1$, $1 \leq k \leq n_2$. We obtain

$$|d_{G \lor H}(u_i, v_k) - d_{G \lor H}(u_j, v_l)| = |n_2 d_G(u_i) + n_1 d_H(v_k) - d_G(u_i) d_H(v_k) - (n_2 d_G(u_i) + n_1 d_H(v_l) - d_G(u_i) d_H(v_l))|$$

Further, by simple algebraic manipulation and by the triangle inequality, we have

$$|d_{G\vee H}(u_{i}, v_{k}) - d_{G\vee H}(u_{j}, v_{l})| \leq n_{2} |d_{G}(u_{i}) - d_{G}(u_{j})| + n_{1} |d_{H}(v_{k}) - d_{H}(v_{l})| + \frac{1}{2} (d_{G}(u_{i}) + d_{G}(u_{j})) |d_{H}(v_{k}) - d_{H}(v_{l})| + \frac{1}{2} (d_{H}(v_{k}) + d_{H}(v_{l})) |d_{G}(u_{i}) - d_{G}(u_{j})|.$$

$$(14)$$

From (13) and (14), we obtain

$$\operatorname{irr}_{t}(G \vee H) \leq \frac{1}{2} \sum_{\substack{u_{i}, u_{j} \in V(G), \\ v_{k}, v_{l} \in V(H)}} [n_{2} |d_{G}(u_{i}) - d_{G}(u_{j})| + n_{1} |d_{H}(v_{l}) - d_{H}(v_{k})|]$$

$$+ \frac{1}{4} \sum_{\substack{u_{i}, u_{j} \in V(G), \\ v_{k}, v_{l} \in V(H)}} (d_{G}(u_{i}) + d_{G}(u_{j})) |d_{H}(v_{k}) - d_{H}(v_{l})|$$

$$+ \frac{1}{4} \sum_{\substack{u_{i}, u_{j} \in V(G), \\ v_{k}, v_{l} \in V(H)}} (d_{H}(v_{k}) + d_{H}(v_{l})) |d_{G}(u_{i}) - d_{G}(u_{j})|.$$

$$(15)$$

The first sum in (15) is equal to $n_2^3 \operatorname{irr}_t(G) + n_1^3 \operatorname{irr}_t(H)$, the second to $2n_1m_1 \operatorname{irr}_t(H)$, and the third to $2n_2m_2 \operatorname{irr}_t(G)$. Hence,

$$\operatorname{irr}_{t}(G \vee H) \leq n_{2}^{3} \operatorname{irr}_{t}(G) + n_{1}^{3} \operatorname{irr}_{t}(H) + 2 n_{1} m_{1} \operatorname{irr}_{t}(H) + 2 n_{2} m_{2} \operatorname{irr}_{t}(G)$$

= $n_{2} (n_{2}^{2} + 2m_{2}) \operatorname{irr}_{t}(G) + n_{1} (n_{1}^{2} + 2m_{1}) \operatorname{irr}_{t}(H).$

2.8 Symmetric difference

The symmetric difference $G \oplus H$ of simple undirected graphs G and H with $|V(G)| = n_1$ and $|V(H)| = n_2$ is the graph with the vertex set $V(G \oplus H) = V(G) \times V(H)$ and the edge set $E(G \oplus H) = \{(u_i, v_k)(u_j, v_l) : \text{ either } u_i u_j \in E(G) \text{ or } v_k v_l \in E(H)\}$. It holds that $|V(G \oplus H)| = n_1 n_2$, and $d_{(G \oplus H)}(u_i, v_j) = n_2 d_G(u_i) + n_1 d_H(v_j) - 2d_G(u_i) d_H(v_j)$ for all $1 \leq i \leq n_1, 1 \leq j \leq n_2$.

Much as in the previous case, we present only the bound on the total irregularity of symmetric difference of two graphs.

Theorem 2.8. Let G and H be simple undirected graphs with $|V(G)| = n_1$, $|E(G)| = m_1$, $|V(H)| = n_2$ and $|E(H)| = m_2$. Then,

$$\operatorname{irr}_t(G \oplus H) \le n_2 (n_2^2 + 4m_2) \operatorname{irr}_t(G) + n_1 (n_1^2 + 4m_1) \operatorname{irr}_t(H).$$

3 Conclusion

In this paper we consider the total irregularity of simple undirected graphs under several graph operations. We present sharp upper bounds for join, lexicographic product, Cartesian product, strong product, direct product and corona product. It is an open problem if the presented upper bounds on the total irregularity of disjunction and symmetric difference are the best possible.

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